A Framework for Control System Design Subject to Average Data-rate Limits

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Networked Control Systems (NCSs)



- NCSs are control systems with communication constraints.
- Key questions in NCSs relate to how channel/network artifacts affect control objectives.
- Typical artifacts include data-rate limits (quantization), random time delays and data dropouts.

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Background: Can information theory help us?

- "Straight out of the box" information theory does not help: Most information theoretic results do not take causality, stability or delays into account, and are asymptotic in nature.
- Nonetheless, there has been advances in NCS subject to average data-rate constraints.

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 Key result (Data-Rate Theorem; Nair and Evans, 2004): In a stochastic setting, and if R is the average data-rate at which data is transmitted over the feedback path,

$$\Re > \sum_{i=1}^{n_p} \ln |p_i|$$

is necessary and sufficient for Mean Square Stability (MSS).

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is necessary and sufficient for Mean Square Stability (MSS).

• The coding schemes that achieve average data-rates close to the above minimum are complex.

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- Key result (Tatikonda et al, 2004; Nair et al, 2007): For fully observed plants, and quadratic criterion, certainty equivalence and (quasi-)separation holds.
- However, no explicit characterization is provided in Nair et al, 2007.
- For fully observed Gaussian first order autoregressive systems, Tatikonda et al, 2004, provides an expression for a causal rate-distortion function. It is not clear if the results are operationally tight or not.

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In this talk we will:

• Present a framework to deal with average data-rate constrained control systems.

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In this talk we will:

- Present a framework to deal with average data-rate constrained control systems.
- Derive achievable bounds on the average data rate needed to achieve a certain performance level for a specific NCS.

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1 Problem Setup

- 2 A lower bound on average data-rates
- 3 Independent source coding schemes
- Interplay between average data-rates and MSS/performance

5 Conclusions

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Problem Setup: Assumptions



- *P* is LTI and has been designed assuming u = y.
- Open loop transfer from *u* to *y* is SISO and strictly proper.
- Initial state and disturbance d are jointly Gaussian, second order; d is stationary with PSD S_d = Ω_dΩ^H_d.

Problem Setup: Assumptions



- *P* is LTI and has been designed assuming u = y.
- Open loop transfer from *u* to *y* is SISO and strictly proper.
- Initial state and disturbance d are jointly Gaussian, second order; d is stationary with PSD $S_d = \Omega_d \Omega_d^H$.
- The channel is a noiseless digital channel with countable alphabet.
- To be able to use the channel, proper encoding and decoding is needed.



• Study the interplay between the average rate at which s_c is generated and the performance and stability of the NCS.

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- Study the interplay between the average rate at which s_c is generated and the performance and stability of the NCS.
- By stability we mean MSS, i.e.,

$$\lim_{k\to\infty} \mathcal{E}\left\{\bar{x}(k)\bar{x}(k)^{T}\right\} < \infty,$$

where \bar{x} is the state of the dynamic system in the figure.



- Study the interplay between the average rate at which *s_c* is generated and the performance and stability of the NCS.
- As performance measure we use the stationary variance of e, namely σ_e^2 .



- Study the interplay between the average rate at which s_c is generated and the performance and stability of the NCS.
- Average data rate refers to

$$\mathcal{R} \triangleq \lim_{k \to \infty} \frac{1}{k} \sum_{i=0}^{k-1} R(i),$$

where R(i) is the expected length (in, e.g., bits or nats) of the symbol sent at time instant *i*, i.e., of $s_c(i)$.

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2 A lower bound on average data-rates

Independent source coding schemes

Interplay between average data-rates and MSS/performance

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A lower bound on average data-rates



• We focus on causal encoders and decoders:

$$s_c(k) = \mathcal{E}_k(y^k, s_c^{k-1}, S_{\mathcal{E}}^k), \quad u(k) = \mathcal{D}_k(s_c^k, S_{\mathcal{D}}^k),$$

where $S_*(k)$ is side info available at * at time k, and $x^k \triangleq x(0), \cdots, x(k)$.

• If $S_{\mathcal{D}} \perp (d, x_o)$ and, upon knowledge of $S_{\mathcal{D}}^k$, \mathcal{D}_k is invertible, then

$$\mathcal{R} \geq I_{\infty}(y \rightarrow u) \triangleq \lim_{k \rightarrow \infty} \frac{1}{k} \sum_{i=0}^{k-1} I(u(i); y^i | u^{i-1}).$$

where I(a, b|c) denotes the conditional mutual information.

 I_∞(y → u) is the directed mutual information rate between y and u (Massey, 1990).

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Independent source coding schemes

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Independent source coding schemes: Definition



A source coding scheme is independent iff:

• $S_{\mathcal{D}} \perp (d, x_o)$, \mathcal{D}_k is invertible given $S_{\mathcal{D}}^k$, and

• The noise
$$n \triangleq u - y$$
 satisfies

$$n=\Omega\cdot q,$$

where q is an i.i.d. zero mean second order sequence, $q \perp (d, x_o)$, and $\Omega \in \mathcal{U}_{\infty}$ (stable, minimum phase and biproper) has deterministic initial state.

Independent source coding schemes: Rewriting

• Equivalent rewriting:



where $A \in \mathcal{U}_{\infty}$, $F \in \mathcal{RH}_2$ (stable strictly proper).

Independent source coding schemes: Rewriting

• Equivalent rewriting:



where $A \in \mathcal{U}_{\infty}$, $F \in \mathcal{RH}_2$ (stable strictly proper).

• In addition, $I_{\infty}(y \rightarrow u) = I_{\infty}(v \rightarrow w)$.

Characterizing $I_\infty(v o w)$ is not easy. We next provide bounds

In an independent source coding scheme,

$$\frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_w}{\sigma_q^2}d\omega \leq I_{\infty}(v \to w) \leq \frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_w}{\sigma_q^2}d\omega + D(q||q_G),$$

where S_w is the stationary PSD of w and σ_q^2 is the variance of q, and $D(q, q_G)$ is the Kullback-Leibler distance between the distribution of q(k) and its Gaussian counterpart.

• These are simple bounds on the directed mutual information rate across an independent source coding scheme.

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Independent source coding schemes: nice, but...



• Q: How to implement the noise q in an independent source coding scheme?

Independent source coding schemes: nice, but...



- Q: How to implement the noise q in an independent source coding scheme?
- A: use an entropy coded dithered quantizer (ECDQ; Zamir and Feder, 1992).

Independent source coding schemes: ECDQs



Q is a scalar uniform quantizer

$$Q(x)=i\Delta, \quad \Delta\left(i-rac{1}{2}
ight)\leq x<\Delta\left(i+rac{1}{2}
ight), i\in\mathbb{Z}.$$

• The dither d_h is i.i.d., $d_h \perp (d, x_o)$, and $d_h(k) \sim \mathcal{U}\left(-\frac{\Delta}{2}, \frac{\Delta}{2}\right)$.

Independent source coding schemes: ECDQs



- EC and ED constitute a loss-less entropy coding pair that work conditioned upon (d^k_h, s^{k-1}).
- Since EC and ED are loss-less, the "outside world" is not affected by their choice.

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Independent source coding schemes: ECDQs



- For any EC and ED, even if there exits strictly causal feedback from *w* to *v* (extension of Zamir and Feder 1992, 1995):
 - q ≜ w − v is i.i.d.,independent of (d, x_o), and second order provided Δ < ∞.

•
$$q(k) \sim \mathcal{U}\left(-\frac{\Delta}{2}, \frac{\Delta}{2}\right)$$
. Thus, $\sigma_q^2 = \Delta^2/12$.

Independent source coding schemes: ach. avg. data-rates



It can be shown that

$$I_{\infty}(v
ightarrow w) \leq \Re < I_{\infty}(v
ightarrow w) + \ln 2,$$

where the gap originates in the inefficiency of the EC (extension of Zamir et al, 2008).

Independent source coding schemes: ach. avg. data-rates



Thus,

$$\frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_w}{\sigma_q^2}d\omega \leq \Re < \frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_w}{\sigma_q^2}d\omega + \frac{1}{2}\ln\left(\frac{2\pi e}{12}\right) + \ln 2,$$

where the second term is the divergence of a uniform distribution from Gaussianity, and the third one is the inefficiency of the EC.

Independent source coding schemes: Whitening w



Since an independent source coding scheme has sufficient degrees of freedom,

- One can always assume, w.l.o.g., that w is white, which also implies that
- the EC-ED pair need only work conditioned upon d_h(k) (memoryless ECDQ).

Independent source coding schemes: recap



For any independent source coding scheme that uses an ECDQ,

$$\frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_{\mathsf{w}}}{\sigma_q^2}d\omega \leq \Re < \frac{1}{4\pi}\int_{-\pi}^{\pi}\ln\frac{S_{\mathsf{w}}}{\sigma_q^2}d\omega + \frac{1}{2}\ln\left(\frac{2\pi e}{12}\right) + \ln 2,$$

Independent source coding schemes: recap



For any independent source coding scheme that uses an ECDQ,

$$\begin{split} \frac{1}{2}\ln\left(1+\gamma\right) &\leq \mathcal{R} < \frac{1}{2}\ln\left(1+\gamma\right) + \frac{1}{2}\ln\left(\frac{2\pi e}{12}\right) + \ln 2,\\ \text{where} \qquad \gamma &\triangleq \frac{\sigma_v^2}{\sigma_q^2} \end{split}$$

is the SNR of the independent source coding scheme.

Independent source coding schemes: recap



For any independent source coding scheme that uses an ECDQ,

$$\frac{1}{2}\ln\left(1+\gamma\right) \leq \mathcal{R} < \frac{1}{2}\ln\left(1+\gamma\right) + \frac{1}{2}\ln\left(\frac{2\pi \mathsf{e}}{12}\right) + \ln 2.$$

 Thus, if we study the interplay between γ and performance (or MSS), then we will be able to say something about the interplay between R and performance (or MSS).



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Average data-rates and MSS:

• Minimal SNR compatible with MSS (Braslavsky et al, 2007; Silva et al, 2010):

$$\inf_{\substack{A,F,\sigma_q^2\\MSS}} \gamma = \prod_{i=1}^{n_p} |p_i|^2 - 1.$$

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Thus, the minimal avg. data-rate compatible with stability satisfies

$$\sum_{i=1}^{n_p} \ln |p_i| \leq \inf_{\substack{A,F,\Delta\\EC-ED\\MSS}} \mathcal{R} < \sum_{i=1}^{n_p} \ln |p_i| + \underbrace{\frac{1}{2} \ln \left(\frac{2\pi e}{12}\right) + \ln 2}_{1.254 \text{bits}}.$$

Just 1.254 bits/sample away from absolute minimal average data-rate for MSS.

We first consider the problem of finding, for $D \in \mathbb{R}^+$,

$$\gamma_{D} \triangleq \inf_{\substack{A, F, \sigma_{q}^{2} \\ MSS, \ \sigma_{e}^{2} < D}} \gamma.$$

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Average data-rates and performance: Prelude II

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If
$$||T_{de}\Omega_d||_2^2 < D < \infty$$
, then
 $\gamma_D = h(\lambda_D) \triangleq \exp\left(\frac{1}{\pi} \int_{-\pi}^{\pi} \ln\left(\sqrt{\frac{Y^2}{\lambda_D} + |S|^2} + \frac{Y}{\sqrt{\lambda_D}}\right) d\omega\right) - 1,$

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Average data-rates and performance: Prelude II

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where $\lambda_D \in \mathbb{R}^+$ is such that

$$g(\lambda_D) \triangleq ||T_{de}\Omega_d||_2^2 + \frac{1}{2\pi} \int_{-\pi}^{\pi} \frac{\lambda_D Y}{2\left(\sqrt{Y^2 + \lambda_D |S|^2} + Y\right)} d\omega = D,$$

and Y, S, T_{de} are given functions of P. (Finding λ_D is easy because g is convex and monotone)

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Average data-rates and performance: main result

• Now, we will focus on average data-rates:

$$\mathcal{R}_{D} \triangleq \inf_{\substack{A,F,\Delta\\EC-ED\\MSS, \ \sigma_{e}^{2} \leq D}} \mathcal{R}.$$

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Average data-rates and performance: main result

Now, we will focus on average data-rates:

$$\mathcal{R}_{D} \triangleq \inf_{\substack{A,F,\Delta\\ EC-ED\\MSS, \ \sigma_{e}^{2} \leq D}} \mathcal{R}.$$

• If
$$||\mathcal{T}_{de}\Omega_d||_2^2 < D < \infty$$
, then
 $\frac{1}{2}\ln(1+\gamma_D) \le \Re_D < \frac{1}{2}\ln(1+\gamma_D) + \frac{1}{2}\ln\left(\frac{2\pi e}{12}\right) + \ln 2$,

where γ_D is as before.

Average data-rates and performance: remarks

• To the best of our knowledge, our results are the first results that, in closed form, provide insights into the interplay between achievable average data-rates and performance in NCS.

Average data-rates and performance: remarks

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- If (x_o, d) is not Gaussian, then the upper bounds on R still hold, but it is difficult to characterize lower bounds.

Average data-rates and performance: remarks

- To the best of our knowledge, our results are the first results that, in closed form, provide insights into the interplay between achievable average data-rates and performance in NCS.
- If (x_o, d) is not Gaussian, then the upper bounds on R still hold, but it is difficult to characterize lower bounds.
- Warning: our bounds may be very conservative when compared with the, still unknown, minimal values of \mathcal{R} that are achievable when causal but otherwise unconstrained source coding schemes are employed.

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- We established a (restricted) bridge between control and information theories using the notion of independent source coding scheme.
- We used that bridge to analyze the interplay between average data-rates and MSS and performance of a specific NCS.
- A key aspect of our approach is that the internal coding scheme SNR plays a fundamental role.
- Open problems:
 - Consider causal but otherwise unconstrained source coding schemes.
 - Consider general control architectures (partly solved).

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- E.I. Silva, M.S. Derpich, J. Ostergaard. A framework for control system design subject to average data-rate constraints. Conditionally accepted in IEEE Transactions on Automatic Control, March 2010.
- E.I. Silva, M.S. Derpich, J. Ostergaard. *Best LTI Performance under Average Data-rate Constraints*. Submitted to the IFAC Workshop on Distributed Estimation and Control in Networked Systems, 2010.
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Thanks for your attention.

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